

Reliable LoadExecs



Topics Covered

- ▶ LoadExec Explained
 - What is LoadExec
 - Advantages
 - Uses
 - Internally
- ▶ Disadvantages of LoadExec
- ▶ Alternatives to LoadExec
 - Overlays
 - Load then Exec
- ▶ Q&A

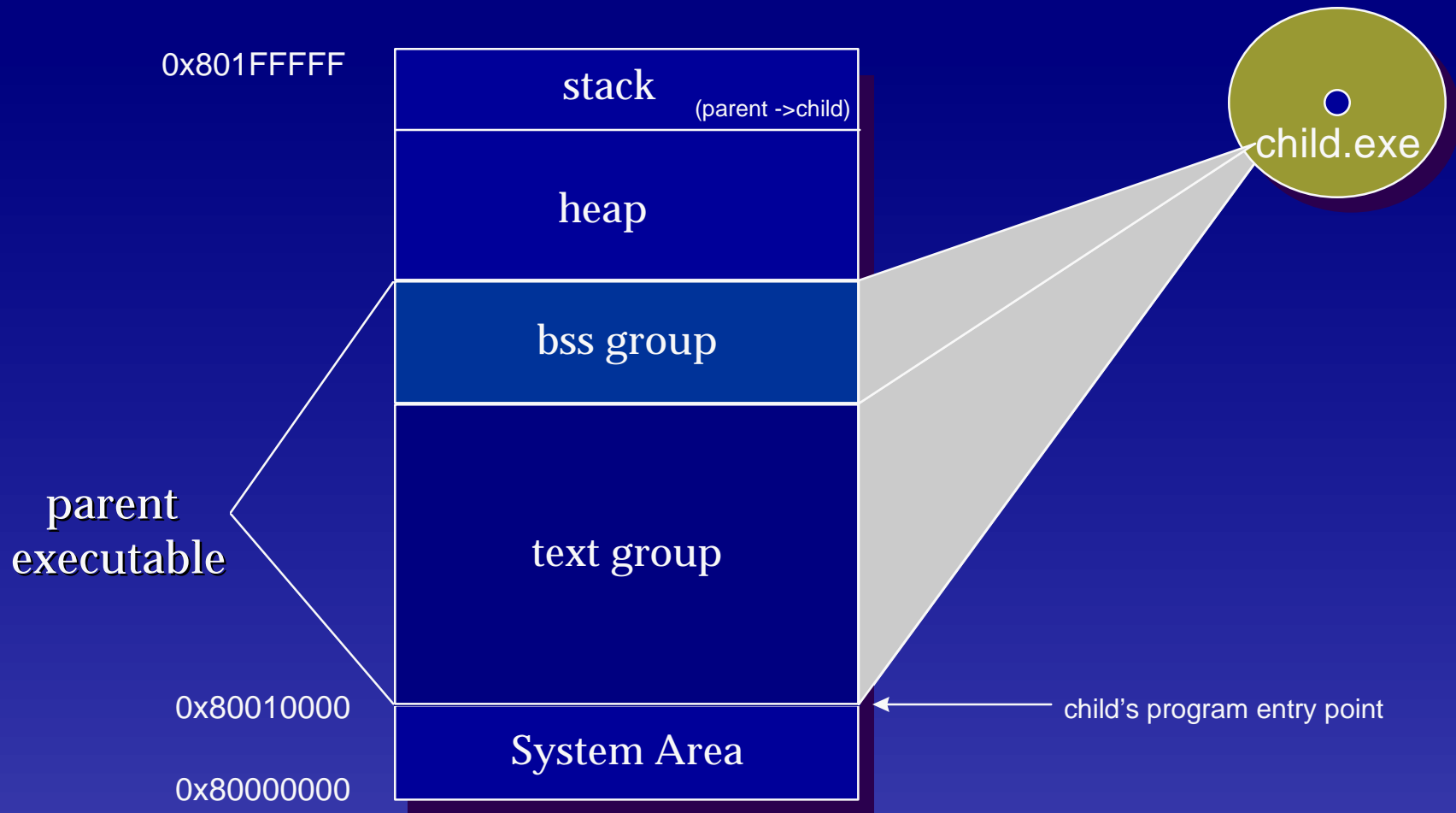
LoadExec Explained

What is Load Exec?

- ▶ LoadExec is a Libapi function that loads an executable from CD into main memory and jumps to the program entry.
 - Calls Load and Exec Internally
 - Executes from ROM
 - Loads child over parent

```
long LoadExec( name, s_addr, s_size )  
char *name; /*executable name*/  
unsigned long s_addr, s_size; /*stack vals*/
```

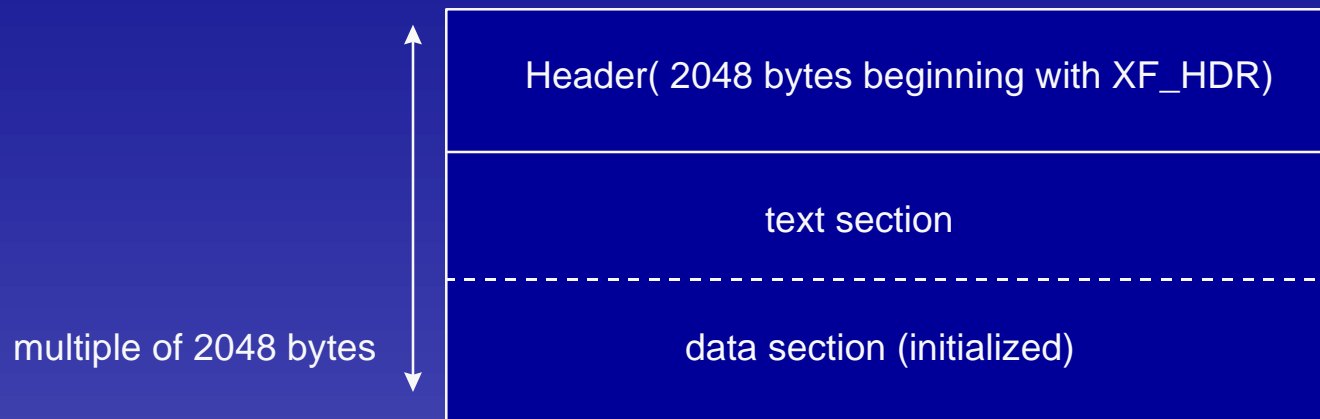
What is LoadExec?



What is LoadExec?

▶ PSX Executable format (as defined in kernel.h)

```
struct XF_HDR {  
    char key[8];           /* key code*/  
    unsigned long text;   /* size of text section*/  
    unsigned long data;   /* size of data section*/  
    struct EXEC exec;     /* executable info*/  
    char title[60];       /* license code */  
};
```



What is LoadExec?

► PSX Executable format (as defined in kernel.h)

```
struct EXEC {  
    unsigned long pc0;           /*initial value of the program counter*/  
    unsigned long gp0;           /*global pointer*/  
    unsigned long t_addr;        /* start address of the text group*/  
    unsigned long t_size;        /*size of the text group*/  
    unsigned long d_addr;        /*system reserved*/  
    unsigned long d_size;        /*system reserved*/  
    unsigned long b_addr;        /*start of the bss group*/  
    unsigned long b_size;        /*size of the bss group*/  
    unsigned long s_addr;        /*lead address of the stack*/  
    unsigned long s_size;        /*stack size*/  
    unsigned long sp,fp,gp,ret,base; /*register shunt area*/  
};
```

Advantages

- ▶ Allows applications to overcome the 2 MB barrier
- ▶ Child processes may use all of Memory
- ▶ Allows parallel development
- ▶ Easy to implement

LoadExecs Uses

- ▶ Demos/Samplers
- ▶ Multi-disc titles
- ▶ To move between exclusive segments of a game
 - Intro -> game
 - Levels
 - Cinematic sequences

LoadExec Internally

- ▶ Saves the current stack values
- ▶ Calls `ExitCriticalSection()`
- ▶ Calls `Load()`

```
addr = exec.s_addr;  
size = exec.s_size;  
ExitCriticalSection();  
if(Load(fn,&exec)==1)
```

LoadExec Internally

▶ Load()

- Fills in system maintained EXEC struct *exec*, with the requested executable's information.
- Loads text group to *exec->t_addr*
- Calls FlushCache()
 - Bad to do outside of a critical section

LoadExec Internally

- ▶ If Load() fails...
 - Attempts to Load() default executable
- ▶ If Load() is successful...
- ▶ Calls Exec()
 - Passing the address of the system EXEC *exec* as an argument

LoadExec Internally

▶ Exec()

- Saves *sp*, *gp*, *fp*, *ra* registers to shunt area in *exec*
- Clears bss
 - If *exec->b_size* $\neq 0$, the programs bss group is cleared using *exec->b_addr* and *exec->b_size*
- Sets-up stack
 - If *exec->s_addr* $\neq 0$, set *sp* and *fp* to *exec->s_addr + exec->s_size*
- Shifts *a2->a1->a0* to allow retrieval of arguments
- Pushes *exec->pc* to a temp reg and does a *ja* on that reg

Disadvantages of LoadExec

Disadvantages of LoadExec

- ▶ When invoking:
 - All interrupts must be stopped
 - Must call `_96_init()`
 - slow
- ▶ Difficult to pass data between processes

Disadvantages of LoadExec

- ▶ Once called all control is relinquished
 - No error handling
 - Load() must successful on first try
 - read() done within in Load also only gets one shot
 - No returning to parent
 - Not because parent is overwritten, but...

Disadvantages of LoadExec

- ▶ If Load() fails on the requested executable...
and...
Load() fails on the default executable...
- ▶ LoadExec enters an infinite loop!!!
 - while(1) { }

Disadvantages of LoadExec

- ▶ The Moral of the Story is...

There is no 100% reliable LoadExec()

Alternatives to LoadExec

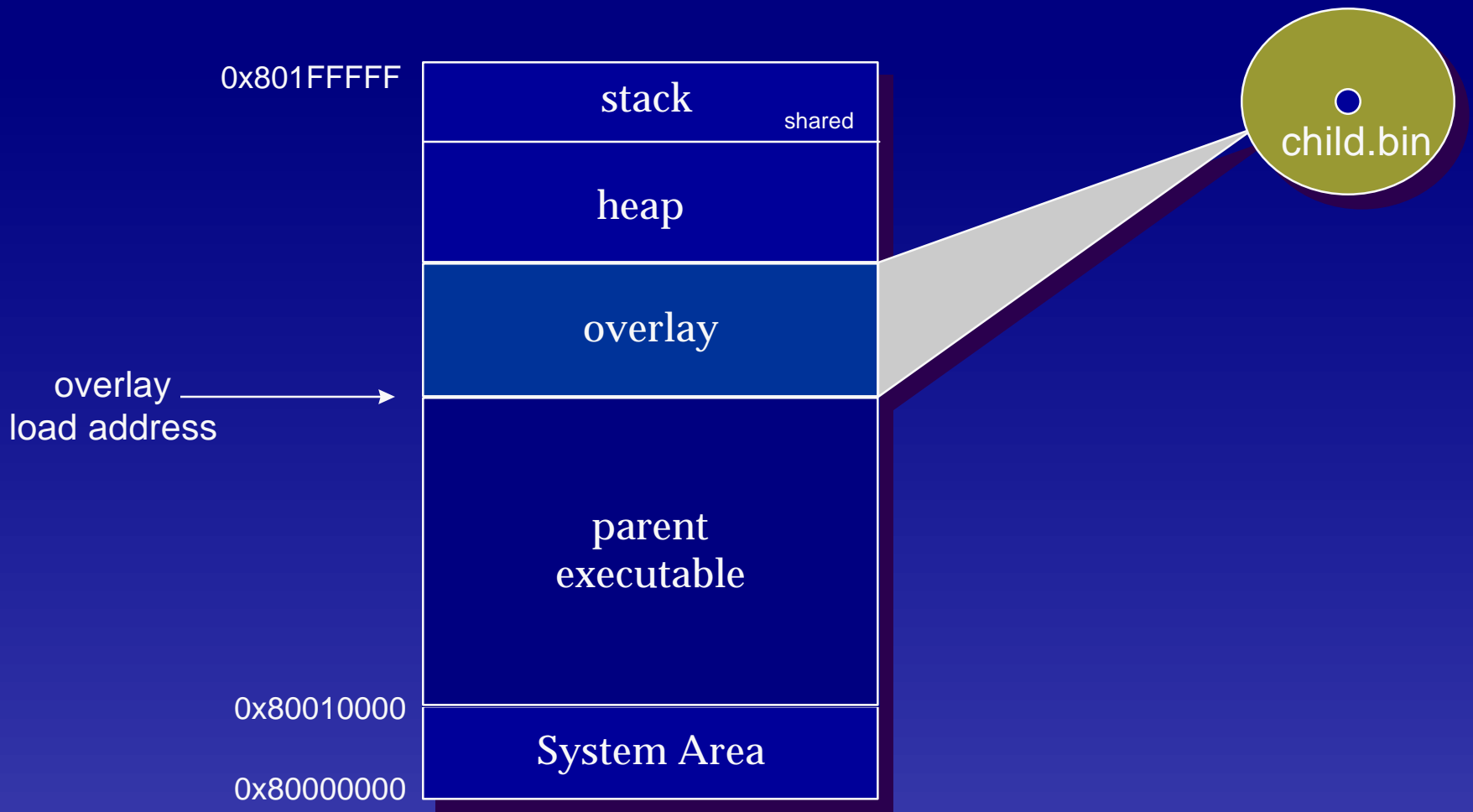
Alternatives to LoadExec

- ▶ Overlays
- ▶ Load then Exec

Overlays

- ▶ A method of linking portions of your code and data to the same memory location then swapping them in and out as needed.

Overlays



Overlays

▶ Advantages:

- Not necessary to stop interrupts.
- A separate process can be executed while the child is loading.
- Can pass arguments and share data.

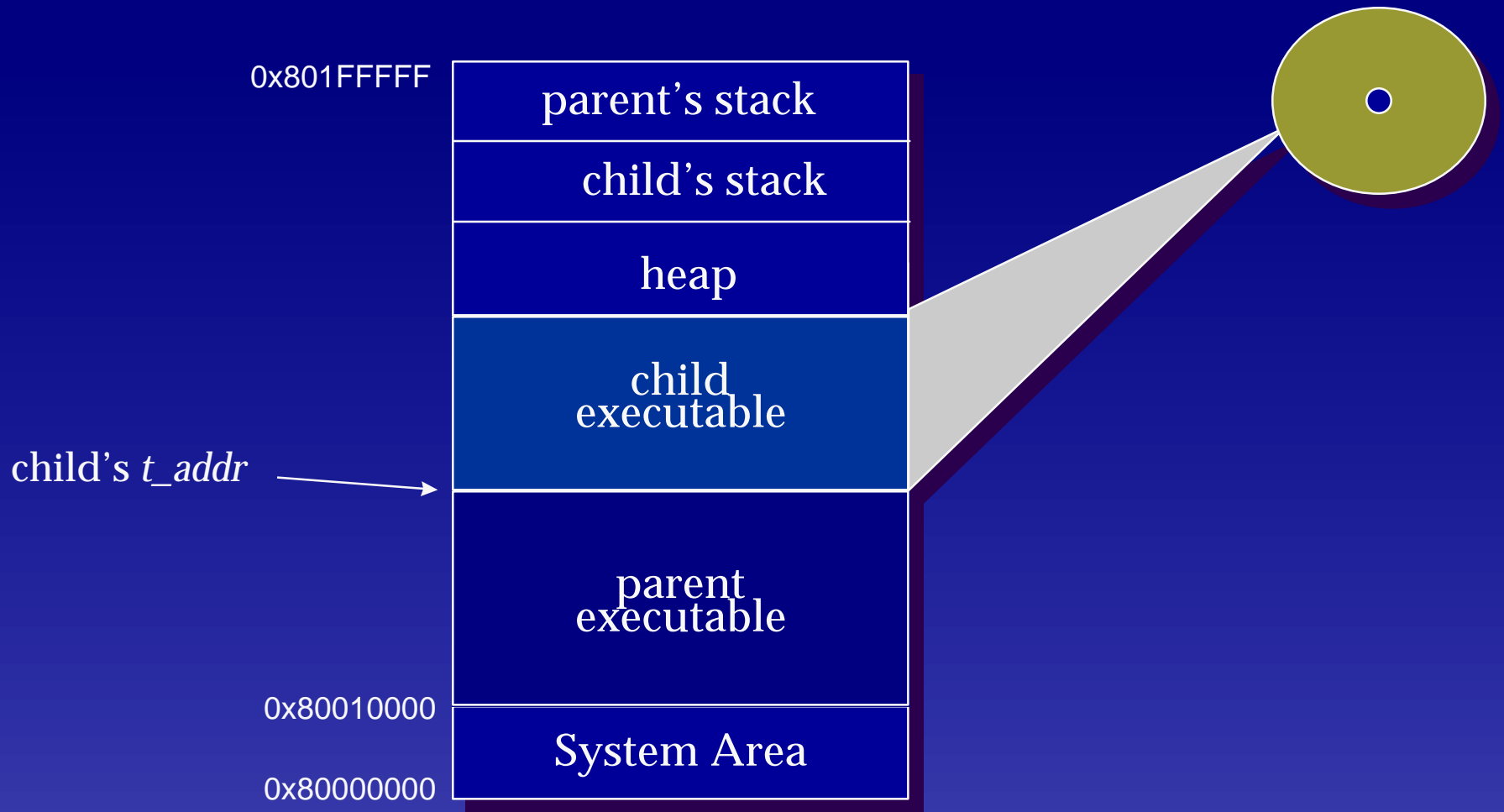
▶ Disadvantages:

- The make operation is more complex.
- The two processes must share memory.
- Processes must be exclusive.

Load then Exec

- ▶ Application calls Load() to get executable from CD into main memory, then calls Exec() to pass control to that executable.

Load then Exec



Load then Exec

▶ Advantages:

- Allows for parallel development
- More control
 - Allows for error handling
 - Application can use own loading scheme i.e. CdRead
- Return to parent is possible
- Easier data sharing

▶ Disadvantages:

- Interrupts must be stopped
- Two copies of Libraries
- Parent and child must operate in limited memory

Load then Exec

But I need more memory!!!!



- ▶ Implement loader scheme
 - Small piece of code that exist only to load and launch larger modules.

Load then Exec - Loader

▶ Loader

- small
- linked just after System Area or in other "safe" region

▶ Child modules

- Link just beyond loader
- Link with none2.obj and explicitly set stack below Loader's stack

Tips...

▶ Parent exit procedures

- ResetGraph(0) or (3)
- CardStop()
- PadStop()
- StopCallback()
- CloseEvent()
- _96_init
- LoadExec() or Load()

Tips...

▶ If Load()

- `_96_remove`
- `EnterCriticalSection()`
- `FlushCache()`
- `Exec()`

▶ Child start up procedures

- `ResetCallback()`
- `CdInit()`
- `InitPAD()`

Tips...

- ▶ If using Load then Exec
 - Repeat child start-up procedures upon return to the parent
- ▶ LoadTest(char *name, EXEC exec)
 - Loads EXEC portion of executable header
 - Returns 0 in case of failure
 - Call prior to Load() or LoadExec()
 - retry if failure
 - if retries unsuccessful inform user to change or clean disc

Tips...

- ▶ Burn CD's at single speed

Q & A

The End